

# Quantification

## 1 Quantifiers

The DPs we have studied so far have generally been of type  $e$ . Let's now consider subject DPs like *everyone*, *no one*,<sup>1</sup> and *someone*.

- (1) Everyone sleeps.

Option 1: Include "plurals" in  $D_e$ , including a symbol that refers to 'nothing,'  $\epsilon$ . *Everyone* is type  $e$ , the sum of all individuals.

- (2) a.  $D_e = \left\{ \begin{array}{l} \epsilon, \text{ Alex, Brie, Cara,} \\ \text{Alex + Brie, Alex + Cara, Brie + Cara,} \\ \text{Alex + Brie + Cara} \end{array} \right\}$
- b.  $\llbracket \text{everyone} \rrbracket = \text{Alex + Brie + Cara (type } e)$
- c.  $\llbracket \text{everyone sleeps} \rrbracket = \text{Sleep(Alex + Brie + Cara)}$

This sort of works for *everyone*, but it does not work for *no one* and *someone*. Why?

Option 2: *Everyone* is not type  $e$ .

- (3) a.  $\llbracket \text{everyone} \rrbracket = \lambda Q_{\langle e, t \rangle} . \forall x [\text{Animate}(x) \rightarrow Q(x)]$
- b.  $\llbracket \text{everyone sleeps} \rrbracket = \forall x [\text{Animate}(x) \rightarrow \text{Sleep}(x)]$

Quantificational DPs are type  $\langle \langle e, t \rangle, t \rangle$ . In other words, they take the VP as their argument.

### Exercise

- (4) **Every** dog sleeps.

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<sup>1</sup>Although we spell this as two words, "no one," we will treat it as one word, just like *nothing*.

## 2 Determiner meanings

We previously wrote meanings for quantificational determiners as relations between sets:

(5) **Quantificational determiners as set-relations:**

- a.  $every/all(A)(B) = 1$       iff  $A \subseteq B$
- b.  $a/some(A)(B) = 1$       iff  $A \cap B \neq \emptyset$
- c.  $no(A)(B) = 1$       iff  $A \cap B = \emptyset$
- d.  $two(A)(B) = 1$       iff  $|A \cap B| \geq 2$
- e.  $more-than-two(A)(B) = 1$       iff  $|A \cap B| > 2$
- f.  $most(A)(B) = 1$       iff  $|A \cap B| > |A \setminus B|$

Because we normally work with truth conditions and functions, not sets, we have to translate (5a) into non-set terms:

- (6)  $\llbracket \text{every dog sleeps} \rrbracket = 1$  iff  $\{x : \text{Dog}(x)\} \subseteq \{y : \text{Sleep}(y)\}$   
 $\Leftrightarrow \forall z_e \in \{x : \text{Dog}(x)\} [z \in \{y : \text{Sleep}(y)\}]$   
 $\Leftrightarrow \forall z_e [ \underbrace{\text{Dog}(z)}_{\text{every's first argument}} \rightarrow \underbrace{\text{Sleep}(z)}_{\text{every's second argument}} ]$

- (7)  $\llbracket \text{every} \rrbracket = \lambda P_{\langle e,t \rangle} . \lambda Q_{\langle e,t \rangle} . \forall z_e [P(z) \rightarrow Q(z)]$

We can similarly rewrite other quantificational determiners as  $\lambda$  functions of type  $\langle\langle e, t \rangle, \langle\langle e, t \rangle, t \rangle\rangle$ .

- (8)  $\llbracket \text{some} \rrbracket =$

- (9)  $\llbracket \text{no} \rrbracket =$

- (10)  $\llbracket \text{two} \rrbracket =$

For *every*, *a/some*, and *no*, we can write the determiner meanings using predicate logic. For others where we have to refer to the size (cardinality) of sets, we will still have to make reference to sets.

### 3 The definite determiner and presupposition calculation

(11) The black cat is in Texas.

A first approximation:

(12)  $\llbracket \text{the}_{\text{sg}} \rrbracket = \lambda P_{\langle e,t \rangle} . \lambda Q_{\langle e,t \rangle} . \text{there is a unique } x [P(x) \wedge Q(x)]$  <sup>2</sup>

What meaning do we predict for (11)? Is that what (11) means?

- (13) a. The elevator in AS5 is broken.  
 b. The escalator in AS5 is broken.

(14) A “partial” semantics for the definite determiner:<sup>3</sup>

$\llbracket \text{the}_{\text{sg}} \rrbracket = \lambda f : f \in D_{\langle e,t \rangle}$  and there is exactly one  $x$  such that  $f(x) = 1$ .  
 the unique  $y$  such that  $f(y) = 1$

(15)  $\llbracket \text{the black cat} \rrbracket = \text{the unique black cat}$   
 $\rightsquigarrow \underbrace{\text{there exists exactly one black cat}}_{\text{presupposition}}$

(16) **Functional Application (revised; compare to H&K p. 76):**<sup>4</sup>

If  $\alpha$  is a branching node,  $\{\beta, \gamma\}$  is the set of  $\alpha$ 's daughters, then

- $\llbracket \alpha \rrbracket$  is defined if and only if:  $\llbracket \beta \rrbracket$  and  $\llbracket \gamma \rrbracket$  are both defined and  $\llbracket \beta \rrbracket$  is a function whose domain contains  $\llbracket \gamma \rrbracket$ ;
- if defined,  $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket(\llbracket \gamma \rrbracket)$ .

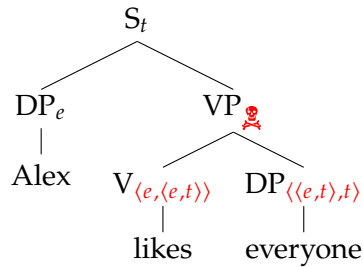
<sup>2</sup>This is not written in valid predicate logic, but we won't dwell on it as we won't adopt this meaning for *the* anyway. If we wanted to, we could write:  $\lambda P_{\langle e,t \rangle} . \lambda Q_{\langle e,t \rangle} . \exists x[P(x) \wedge Q(x) \wedge \neg \exists y[y \neq x \wedge P(y) \wedge Q(y)]]$ .

<sup>3</sup>A *partial function* is a function that is not defined for all possible values of its arguments.

<sup>4</sup>H&K describes this in terms of linguistic objects *being in the domain of*  $\llbracket \cdot \rrbracket$  rather than being defined or not.

## 4 Quantifiers in object position

(17) Alex likes everyone.



We'll first consider the related sentence in (18), and then return to (17).

(18) Everyone, Alex likes \_\_\_\_.

Example (18) involves *movement* (specifically, topicalization) of the object. We need a semantics for how we interpret movement:

(19) **The interpretation of movement:**

Pick an arbitrary variable, such as  $x$ .

a. The base position of movement is replaced with a *trace*;  $\llbracket t \rrbracket = x$ , type  $e$ .

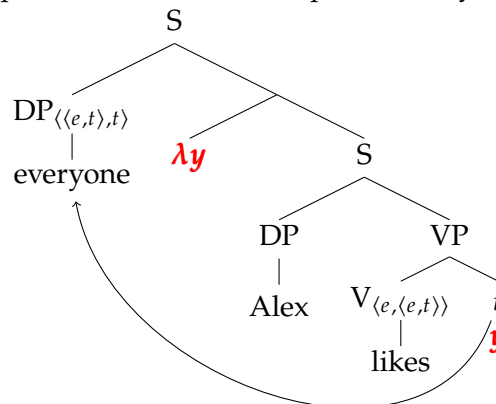
b. A  $\lambda$ -binder  $\lambda x$  is adjoined right under the target position of the movement chain.

(20) **How to interpret  $\lambda$ s in trees (aka  $\lambda$  Rule):**

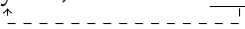
(to be revisited later)

$$\left[ \left[ \lambda x \quad \dots x \quad \dots \right] \right] = \lambda x . \dots x \dots$$

Now notice that objects of type  $\langle \langle e, t \rangle, t \rangle$  can be interpreted easily if they are moved:



Let's now return to example (17). One possible solution to the problem of quantifiers in object position is to interpret the sentence *as if the object has moved*, as in (18). The structure that we interpret is called *Logical Form* (LF).

- (21) a. Surface structure: Alex likes everyone. =(17)  
b. Logical Form (LF): everyone, Alex likes \_\_\_\_.
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Such movement that is not reflected in the surface structure are called *covert*; I use dashed arrows for covert movement. The covert movement of quantifiers as in (21) is called *Quantifier Raising* (QR) (May, 1977). QR is required for quantifiers that are not in subject position, in order to avoid the type problem in (17).

## References

May, Robert Carlen. 1977. The grammar of quantification. Doctoral Dissertation, Massachusetts Institute of Technology.