

Variables, pronouns, and scope

1 Notes on variables

(1) Some math “sentences”:

- a. $1 = 2 - 1$ a sentence with no variables; not context-sensitive
- b. $n = 2 - 1$ a sentence with a variable; context-sensitive
- c. $\forall n (2(n + 1) = 2n + 2)$ a sentence with a variable; *not* context-sensitive

- We say (1b) contains a *free variable* because the truth of the sentence depends on the context. In particular, the sentence is true iff the variable “ n ” is interpreted as 1.
- The truth of sentence (1c), like (1a), does not depend on the context at all.

(2) Some terminology, using (1c) as an example:

$$\underbrace{\forall n}_{\text{binder}} \left(\underbrace{2(n + 1) = 2n + 2}_{\text{scope}} \right)$$

- *Binders* control the interpretation of a particular variable within a certain part of its structure, which we call its *scope*. Here, \forall binds the variable n in its scope.
- We call variables that are in the scope of a matching binder *bound variables*.

Let’s call the mapping between free variables and their values *assignment*.

2 Pronouns

This free/bound terminology is useful for natural language as well:

- (3) a. John likes Mary. a sentence with no variables; not assignment-sensitive
- b. John likes him. a sentence with a variable; assignment-sensitive
- c. Every boy likes himself. a sentence with a variable; *not* assignment-sensitive

We’ll formalize this by giving each pronoun a numerical *index*. We’ll compute denotations relative to an *assignment function*, which is a function from the set of indices (\mathbb{N}) to D_e .

(4) Pronouns Rule (to be replaced later):

If α is a pronoun, g is a variable assignment, and $g(i)$ is defined, then $\llbracket \alpha_i \rrbracket^g = g(i)$.

- (5) Suppose g is a function and $g(3) = \text{Sam} \in D_e$.
 - a. $\llbracket \text{him}_3 \rrbracket^g = \text{Sam}$
 - b. $\llbracket \text{John likes him}_3 \rrbracket^g = 1$ iff John likes Sam

Q: Does it matter what g returns for other values in (5)?

A: No. It might even be undefined for other values.

Q: Why did we use 3? Does the number matter?

A: The choice of number was arbitrary, but it is important whether or not we reuse numbers:

- (6) a. He₂ thinks that he₂ is smart.
b. He₂ thinks that he₇ is smart.

Q: Does the assignment function affect other parts of the sentence?

A: No. “John” and “likes” are *constants*, meaning their values are the same no matter the assignment: for any assignment function f , $\llbracket \text{John} \rrbracket^f = \text{John}$.

Warning: There’s a section of H&K (pp. 92–109) where they just use notation like $\llbracket \text{him} \rrbracket^{\text{John}} = \text{John}$, which only accommodates one variable at a time, but then they introduce their actual notation on page 110, which we use here.

3 Rules with assignments

In order to work with assignment functions, we need to modify all our existing rules so that they pass assignment functions. These definitions are based on H&K p. 95:

(7) **Terminal Nodes (TN):** (unchanged)

If α is a terminal node, $\llbracket \alpha \rrbracket$ is specified in the lexicon.¹

(8) **Non-branching Nodes (NN):**

If α is a non-branching node, and β is its daughter node, then, for any assignment g , $\llbracket \alpha \rrbracket^g = \llbracket \beta \rrbracket^g$.

(9) **Functional Application (FA):**

If α is a branching node, $\{\beta, \gamma\}$ is the set of α ’s daughters, then, for any assignment g , if $\llbracket \beta \rrbracket^g$ is a function whose domain contains $\llbracket \gamma \rrbracket^g$, then $\llbracket \alpha \rrbracket^g = \llbracket \beta \rrbracket^g(\llbracket \gamma \rrbracket^g)$.

(10) **Predicate Modification (PM):**

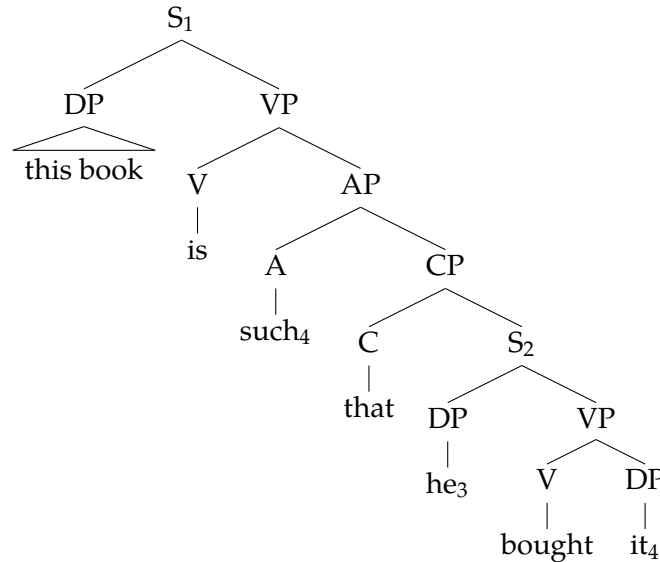
If α is a branching node, $\{\beta, \gamma\}$ is the set of α ’s daughters, then, for any assignment g , if $\llbracket \beta \rrbracket^g$ and $\llbracket \gamma \rrbracket^g$ are both of type $\langle e, t \rangle$, then $\llbracket \alpha \rrbracket^g = \lambda x \in D_e . \llbracket \beta \rrbracket^g(x) = 1$ and $\llbracket \gamma \rrbracket^g = 1$.

¹H&K proposes (p. 94) to still use $\llbracket \alpha \rrbracket$ without an assignment function superscript for *constants*, i.e. if $\llbracket \alpha \rrbracket^g$ is the same value for all assignment functions g .

4 *Such that* relatives

The English expression *such that* allows us to construct relative clauses without movement.²

(11) ? This book is $such_4$ that he_3 bought it_4 . ($g(3) = \text{John}$)



Here, (11) does not seem assignment-dependent. But the Principle of Compositionality states that $\llbracket S_1 \rrbracket$ be computed based on the meaning of $\llbracket S_2 \rrbracket$, which contains a pronoun and is assignment-dependent.

Idea: *Such* binds *it*, doing the work of creating a *predicate* out of the assignment-dependent sentence “John bought it.”

(12) ***Such* Rule (temporary):**³
 $\llbracket such_i \gamma \rrbracket^g = \lambda x_e . \llbracket \gamma \rrbracket^{[i \mapsto x] \parallel g}$

$[i \mapsto x] \parallel g$ is the *combination* of functions $[i \mapsto x]$ and g :

(13) **Definition: function combination**

$$f \parallel g \equiv \lambda x . \begin{cases} f(x) & \text{if } x \in \text{domain}(f) \\ g(x) & \text{otherwise} \end{cases}$$
 Read “ f or else g .”

Let’s compute $\llbracket S_1 \rrbracket^g$ with the following global assignment function: $g = \begin{bmatrix} 3 \mapsto \text{John} \\ 11 \mapsto \text{Tama} \end{bmatrix}$.
 Assume $\llbracket \text{that} \rrbracket = \text{Id}$.

Warning: H&K uses $g^{x/i}$ notation for $[i \mapsto x] \parallel g$, but I think it’s confusing so I don’t use it.⁴

²Unfortunately, the use of *such that* sounds “unlyrical” (Quine, 1960, §23)... but we’ll ignore that here.

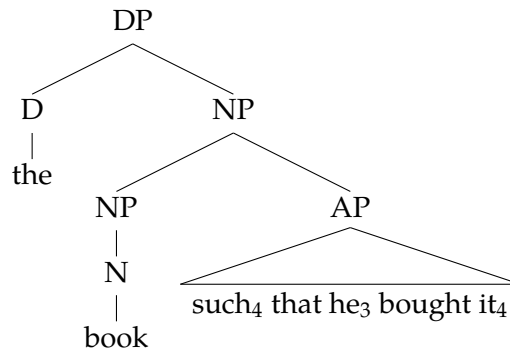
³“Such” does not have a type. That’s why it can only be interpreted using the *Such* Rule.

⁴For one, I’ve also seen very similar notation “ $g(x/a)$ ” for a function that maps x to a , which is the reverse of

We can also use *such that* to construct (slightly awkward) *relative clauses*:

(14) ? the book such_4 that he_3 bought it_4

The semantics for *such* above works perfectly fine here.



Binding multiple variables:

(15) ? This book is such_4 that he_3 bought it_4 and then gave it_4 to Sarah.

(16) ? every book such_4 that he_3 bought it_4 and then gave it_4 to Sarah

Binding no variables (vacuous binding):

(17) * This book is such_4 that today is Monday.

(18) * every book such_4 that today is Monday

The ungrammaticality of these examples shows that binding *no* variables is disallowed by the grammar. This is called *vacuous binding*.

5 Traces & Pronouns

(19) **The interpretation of movement (revised):** replaces last week's movement rule
Pick an arbitrary index i .

a. The base position of movement is replaced with a *trace* with index i : t_i .

b. A *binder index* i is adjoined right under the target position of the movement chain.

(20) **Traces and Pronouns Rule (T&P):** replaces Pronouns Rule in (4)

If α is a pronoun or trace, g is a variable assignment, and $g(i)$ is defined, then

$$\llbracket \alpha_i \rrbracket^g = g(i).$$

(21) **Predicate Abstraction (PA):** (H&K p. 186 version)

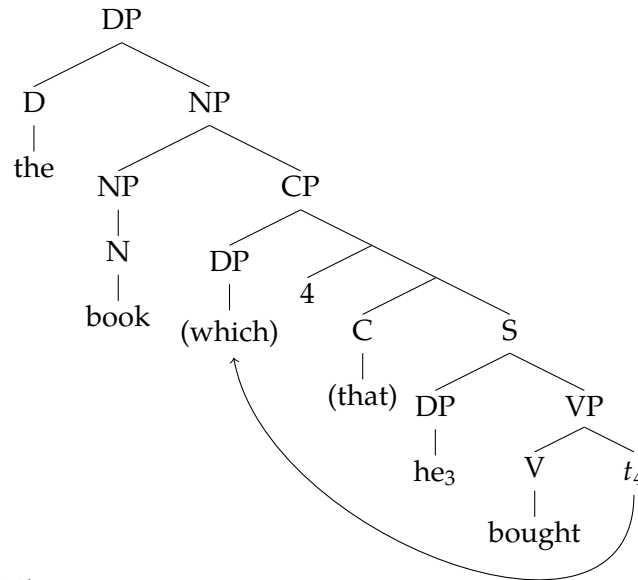
replaces last week's rule for λ nodes in the tree and the *Such* Rule in (12)⁵

Let α be a branching node with daughters β and γ , where β dominates only a numerical index i . Then, for any assignment g , $\llbracket \alpha \rrbracket^g = \lambda x . \llbracket \gamma \rrbracket^{[i \mapsto x] \parallel g}$.

what H&K mean in their x/i .

⁵We can think of "such" as the pronunciation of a lexicalized binder index, not generated through movement.

(22) the book that he₃ bought ___



Exercise: Compute (22).

6 Quantifier scope

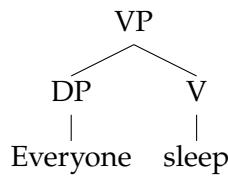
(23) Everyone does not sleep (during class).

- a. $1 \text{ iff } \forall x \in D_e \left[\underbrace{x \text{ is animate} \rightarrow \text{it's not that } \underbrace{[x \text{ sleeps (during class)]}_{\text{scope of not}}}}_{\text{scope of } \forall} \right] \quad (\forall > \text{not})$
- b. $1 \text{ iff it's not that } \left[\underbrace{\forall x \in D_e \left[\underbrace{x \text{ is animate} \rightarrow x \text{ sleeps (during class)}}_{\text{scope of } \forall} \right]}_{\text{scope of not}} \right] \quad (\text{not} > \forall)$

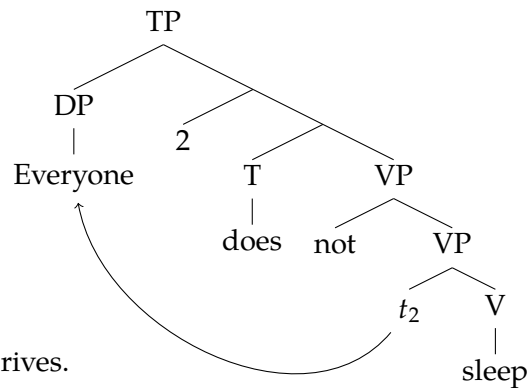
The two readings in (23) represent a *scope ambiguity*. There are two operators that “take scope”— \forall and negation—and one scope contains the other. We say \forall in (23a) takes *wider* scope, and write $\forall > \text{not}$ to indicate this.

Recall from the problem set that there are advantages to adopting a VP-internal subject, interpreted through movement. We will adopt this here.

Step 1: Build subject in Spec,VP



Step 2: Add *not* + T, move subject DP to Spec,TP

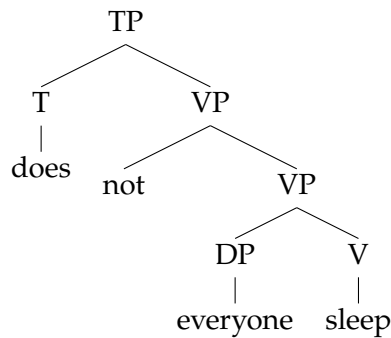


Exercise: Let's see what meaning this tree derives.

We call the meaning that is reflected on the surface form—here, (23a)—a *surface scope* reading.

How do we get reading (23b)? One option: *pretend the movement didn't take place*.

At Logical Form (LF): Pretend the movement didn't happen



Exercise: Interpret this tree.

We call this the *inverse scope* interpretation. The process of “ignoring” movement at LF is called *syntactic reconstruction*.

References

Quine, Willard Van Orman. 1960. *Word and object*. Cambridge.