

# Final review

## 1 Key terms and concepts

Key terms and concepts, roughly in the order they were introduced. (*For terms in italics and parentheses, only a very basic understanding is expected.*)

- truth condition, compositionality
- entailment ( $\Rightarrow$ ), presupposition ( $\rightsquigarrow$ ): how to distinguish
- set notation:  $\{x : x \text{ is a cat}\}$   $\in \subseteq = \cap \cup \setminus$
- other mathematical notation:  $\forall \exists \wedge \vee f \parallel g [i \mapsto x]$
- the interpretation function  $\llbracket \dots \rrbracket$
- types:  $e, t, \langle \sigma, \tau \rangle$   $D_\tau$  is the domain of type  $\tau$
- $\lambda$ -notation:  $f = \lambda \underbrace{x}_{\text{argument variable}} : \underbrace{x \in \mathbb{R}}_{\text{domain condition}} . \underbrace{x + 1}_{\text{value description}}$
- characteristic functions of sets
- The Triangle Method; Binary Branching
- variables: bound vs free, binders, *such that*, index, assignment function, vacuous binding
- The interpretation of movement:  
Pick an arbitrary index  $i$ .
  1. The base position of movement is replaced with a *trace* with index  $i$ :  $t_i$ .
  2. A *binder index*  $i$  is adjoined right under the target position of the movement chain.
- Quantifiers:
  - Generalized Quantifier Theory
  - quantificational determiners
  - type  $\langle \langle e, t \rangle, t \rangle$
  - the problem of quantifiers in object position
  - restrictor, scope, bound pronouns
- Grammatical architecture:
  - T/Y-model: Syntax, Logical Form (LF), Phonological Form (PF)
  - overt and covert movement; islands
  - Quantifier Raising (QR); (*reconstruction*)

- Ellipsis: (Handout 7)
  - LF Identity Condition on Ellipsis (H&K p. 250) (7)
  - scope parallelism
  - strict and sloppy readings of pronouns
  - Antecedent-Contained Deletion (ACD); Sag-Williams generalization (8)
- Focus: (Handout 8)
  - alternatives, preadjacent, F-mark
  - focus-semantic value  $[[\dots]]^f$ , point-wise composition (9), ordinary semantic value  $[[\dots]]^o$
  - *only, also*
  - Taglicht ambiguities
- Modals and conditionals: (Handout 9)
  - extension, intension
  - possible worlds; type  $s$
  - modals bases: epistemic (EPIST), deontic (DEONT), “root”
  - modal forces: possibility ( $\exists$ ), necessity ( $\forall$ )
  - Intensional Functional Application (10)
  - conditionals: material implication ( $\rightarrow$ ), modal restrictor view
- Negative Polarity Items (NPIs): (Handout 10)
  - licenser, upward-entailing, downward-entailing (DE), (*strong NPIs*)
  - *even, (even associating with an indefinite)*
- Questions: (Handout 11)
  - the meaning of a question is..., (*weak vs strong exhaustivity*)
  - question-answer congruence (11)
  - alternative question; multiple *wh*-questions: single-pair vs pair-list, ALTSHIFT (12)
- Binding and coreference: (Handout 12)
  - Traditional Binding Theory: Conditions A, B, C
  - coreference: accidental coreference vs variable binding
  - syntactic binding vs semantic binding
  - Have Local Binding! (13)
  - indices under ellipsis, (*Dahl's puzzle*)

## 2 Basic composition rules

You will have to be able to use these six rules without their definitions, and give their names.

(1) **Terminal Nodes (TN):**

If  $\alpha$  is a terminal node,  $\llbracket \alpha \rrbracket$  is specified in the lexicon.

(2) **Non-branching Nodes (NN):**

If  $\alpha$  is a non-branching node, and  $\beta$  is its daughter node, then  $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket$ .

(3) **Functional Application (FA):** (Handout 4 version; based on H&K)

If  $\alpha$  is a branching node,  $\{\beta, \gamma\}$  is the set of  $\alpha$ 's daughters, then

- $\llbracket \alpha \rrbracket$  is defined if and only if:  $\llbracket \beta \rrbracket$  and  $\llbracket \gamma \rrbracket$  are both defined and  $\llbracket \beta \rrbracket$  is a function whose domain contains  $\llbracket \gamma \rrbracket$ ;
- if defined,  $\llbracket \alpha \rrbracket = \llbracket \beta \rrbracket(\llbracket \gamma \rrbracket)$ .

(4) **Predicate Modification (PM):**

If  $\alpha$  is a branching node,  $\{\beta, \gamma\}$  is the set of  $\alpha$ 's daughters, and  $\llbracket \beta \rrbracket$  and  $\llbracket \gamma \rrbracket$  are both in  $D_{\langle e,t \rangle}$ , then  $\llbracket \alpha \rrbracket = \lambda x \in D_e . \llbracket \beta \rrbracket(x) = 1$  and  $\llbracket \gamma \rrbracket = 1$

(5) **Traces and Pronouns Rule (T&P):**

If  $\alpha$  is a pronoun or trace,  $g$  is a variable assignment, and  $g(i)$  is defined, then  $\llbracket \alpha_i \rrbracket^g = g(i)$ .

(6) **Predicate Abstraction (PA):**

Let  $\alpha$  be a branching node with daughters  $\beta$  and  $\gamma$ , where  $\beta$  dominates only a numerical index  $i$ . Then, for any assignment  $g$ ,  $\llbracket \alpha \rrbracket^g = \lambda x . \llbracket \gamma \rrbracket^{[i \mapsto x]g}$ .

## 3 Additional technical concepts

If any of the following rules/concepts are necessary on the final, they will be given:

(7) **LF Identity Condition on Ellipsis:** (H&K p. 250)

A constituent may be deleted at PF only if it is a copy of another constituent at LF.

A rule for indices (p. 254): No LF representation (for a sentence or multisentential text) must contain both bound occurrences and free occurrences of the same index.

(8) **The Sag-Williams generalization:** (Sag, 1976; Williams, 1974)

In Antecedent-Contained Deletion, the size of the ellipsis determines the lowest possible scope of the object DP.

(9) **A recursive definition for the computation of focus-semantic values:**

Terminal nodes (TN):

$$\llbracket \alpha_\tau \rrbracket^f = \begin{cases} \{\llbracket \alpha_\tau \rrbracket\} & \text{if } \alpha \text{ not F-marked} \\ \text{a subset of } D_\tau & \text{if } \alpha \text{ F-marked} \end{cases}$$

Point-wise functional application (PFA):

$$\left[ \begin{array}{c} \alpha_\tau \\ \swarrow \quad \searrow \\ \beta_{\langle \sigma, \tau \rangle} \quad \gamma_\sigma \end{array} \right]^f = \begin{cases} \{b(g) \mid b \in \llbracket \beta \rrbracket^f, g \in \llbracket \gamma \rrbracket^f\} & \text{if } \alpha \text{ not F-marked} \\ \text{a contextually-determined subset of } D_\tau & \text{if } \alpha \text{ F-marked} \end{cases}$$

(10) **Intensional Functional Application:** (from Handout 9; based on von Stechow and Heim, 2011)

If  $\alpha$  is a branching node and  $\{\beta, \gamma\}$  is the set of its daughters, then, for any world  $w$  and assignment  $g$ : if  $\llbracket \beta \rrbracket^{w,g}$  is a function whose domain contains  $\lambda w'_s . \llbracket \gamma \rrbracket^{w',g}$ , then  $\llbracket \alpha \rrbracket^{w,g} = \llbracket \beta \rrbracket^{w,g} (\lambda w'_s . \llbracket \gamma \rrbracket^{w',g})$ .

(11) **Condition on question-answer congruence (Rooth, 1992):**

$$\llbracket \text{question} \rrbracket^o \subseteq \llbracket \text{answer} \rrbracket^f$$

(12) **ALTSHIFT (Kotek, to appear):**

$$\begin{aligned} \llbracket [\text{ALTSHIFT } \alpha] \rrbracket^o &= \llbracket \alpha \rrbracket^f \\ \llbracket [\text{ALTSHIFT } \alpha] \rrbracket^f &= \left\{ \llbracket [\text{ALTSHIFT } \alpha] \rrbracket^o \right\} = \left\{ \llbracket \alpha \rrbracket^f \right\} \end{aligned}$$

(13) **Have Local Binding! (Büring, 2005):**

For any two NPs  $\alpha$  and  $\beta$ , if  $\alpha$  could semantically bind  $\beta$  (i.e. if it c-commands  $\beta$  and  $\beta$  is not semantically bound in  $\alpha$ 's c-command domain already),  $\alpha$  must semantically bind  $\beta$ , unless that changes the interpretation.

## References

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- Kotek, Hadas. to appear. On the semantics of *wh*-questions. In *Proceedings of Sinn und Bedeutung* 20.
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